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## RESEARCH ARTICLE

### Approximate Invariant Subspaces and Quasi-Newton Optimization Methods

Serge Gratton and Philippe L. Toint

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New approximate secant equations are shown to result from the knowledge of (problem dependent) invariant subspace information, which in turn suggests improvements in quasi-Newton methods for unconstrained minimization. A new limited-memory BFGS using approximate secant equations is then derived and its encouraging behaviour illustrated on a small collection of multilevel optimization examples. The smoothing properties of this algorithm are considered next, and automatic generation of approximate eigenvalue information demonstrated. The use of this information for improving algorithmic performance is finally investigated on the same multilevel examples.

**Keywords:** large-scale optimization, quasi-Newton methods, limited memory algorithms, discretized problems, multilevel optimization.

#### 1. Introduction

The history of quasi-Newton methods for optimization is rich and long. Starting with the Davidon-Fletcher-Powell (DFP) method (Davidon [8], Fletcher and Powell [11]), most famously represented by the Broyden-Fletcher-Goldfarb-Shanno (BFGS) update (Broyden [3], Fletcher [10], Goldfarb [12] and Shanno [28]), and excellently explained in the classical book by Dennis and Schnabel [9], they have played an important role in the solution of practical problems. In the context of unconstrained minimization, i.e. the solution of problems of the form

$$\min_{x \in \mathbb{R}^n} f(x) \quad (1)$$

for a smooth objective function  $f$  from  $\mathbb{R}^n$  into  $\mathbb{R}$ , they attempt to construct, around a given point  $x \in \mathbb{R}^n$ , a second-order model of this function of the form

$$m(x + s) = f(x) + \langle g(x), s \rangle + \frac{1}{2} \langle s, Bs \rangle$$

where  $g(x) \stackrel{\text{def}}{=} \nabla_x f(x)$ , and where  $B$  is an (often positive-definite) approximation of the Hessian matrix  $\nabla_{xx} f(x)$ , capturing information about the curvature of the objective function around  $x$ . Quasi-Newton methods then proceed to exploit a sequence of models of this type in an iterative manner. In this process, the curvature information at iterate  $x_{k+1}$  is obtained by updating the approximate Hessian matrix  $B_k$  to obtain the new approximation  $B_{k+1}$  such that the *secant equation*

$$B_{k+1}s_k = y_k, \quad (2)$$

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holds, where

$$s_k \stackrel{\text{def}}{=} x_{k+1} - x_k \text{ and } y_k \stackrel{\text{def}}{=} g(x_{k+1}) - g(x_k), \quad (3)$$

The pair  $(s_k, y_k)$  is then said to be the quasi-Newton pair associated with equation (2). If positive-definiteness of the matrix  $B_k$  is also maintained throughout the iterations (as can be enforced for instance with the BFGS or DFP updates), the search direction at iteration  $k$  is then computed from

$$d_{k+1} = -B_{k+1}^{-1}g(x_{k+1}) \quad (4)$$

and a linesearch is performed along this direction (see Dennis and Schnabel [9] for details). In order to avoid the cost of solving the linear system in (4), the matrix  $H_{k+1} \stackrel{\text{def}}{=} B_{k+1}^{-1}$  is typically recurred instead of  $B_{k+1}$ , using the *inverse secant equation*

$$H_{k+1}y_k = s_k \quad (5)$$

as an alternative to (2), and

$$d_{k+1} = -H_{k+1}g(x_{k+1}) \quad (6)$$

instead of (4). Note that (5) uses the same pair as (2) but in the reverse order. For the BFGS update, which is particular interest to us in this paper, the relevant updating formula is then given by

$$H_{k+1} = \left( I - \frac{s_k y_k^T}{y_k^T s_k} \right) H_k \left( I - \frac{y_k s_k^T}{y_k^T s_k} \right)^T + \frac{s_k s_k^T}{y_k^T s_k} \quad (7)$$

It readily follows from this formula that  $H_{k+1}$  remains positive-definite if  $H_k$  is positive-definite and if

$$y_k^T s_k > 0, \quad (8)$$

a condition one can always enforce in the linesearch procedure if the objective function is bounded below (again see Dennis and Schnable [9]).

We are especially interested in the application of quasi-Newton to large-scale problems, in which case it is often impractical to store the (dense) matrices  $B_{k+1}$  or  $H_{k+1}$  explicitly. In such a context, a “limited-memory” version of the quasi-Newton method has been pioneered by several authors, in which the matrix  $H_{k+1}$  is assembled at every iteration as a product of finitely many low-rank updates, each involving a pair  $(s_j, y_j)$  (see Liu and Nocedal [22] and Byrd, Lu and Nocedal [5], for the most famous algorithm of this type and further references).

Our purpose in the present paper is to show that additional knowledge about the eigenstructure of the local Hessian matrix  $\nabla_{xx}f$  can be used to advantage in order to capture more information on the local curvature. We discuss in particular how this can be achieved when limited-memory BFGS updates are considered, and illustrate the practical motivation for this analysis in the case of large-scale multilevel unconstrained optimization. Examples of this type are presented, and it is shown that our proposal may improve their numerical solution substantially.

The resulting multilevel optimization method is a linesearch algorithm, at variance with the trust-region based techniques discussed in Gratton, Sartenaer and

Toint [14, 15] or Gratton, Mouffe, Toint and Weber-Mendonça [13]. It also differs from the proposals by Nash [25], Lewis and Nash [20] and Wen and Goldfrab [32] in that none of these techniques makes explicit use of limited-memory quasi-Newton Hessian approximations.

Section 2 introduces the use of invariant subspaces in the derivation of secant information, presents the resulting quasi-Newton algorithm and discusses the specialization of this new algorithm to the multilevel case. Section 4 briefly describes our test problems, and reports our first numerical results. Section 5 discusses further consequences of the use of the new algorithm to multilevel optimization, introduces approximate eigenvalue equations and their use within the new algorithmic framework. Section 6 presents the associated numerical tests. Some conclusions and perspectives are finally outlined in Section 7.

## 2. Invariant Subspaces in Quasi-Newton Methods

### 2.1. Invariant Subspaces and Approximate Secant Equations

We start by considering the case where the objective function  $f$  is a convex quadratic, that is

$$f(x) = f + \langle g, x \rangle + \frac{1}{2} \langle x, Gx \rangle$$

where  $f \in \mathbb{R}$ ,  $g \in \mathbb{R}^n$  and  $G$  is a positive-definite symmetric matrix. In this case, it is easy to verify that, if we consider a step  $s_k$  at iteration  $k$  of a quasi-Newton algorithm, then

$$Hy_k = s_k \tag{9}$$

where  $H = G^{-1}$ . Assume now that we also know a decomposition of  $\mathbb{R}^n$  in a collection of invariant subspaces  $\{\mathcal{S}_i\}_{i=1}^p$  related to  $G$ , i.e. subspaces such that, for each  $i$ ,  $Gd \in \mathcal{S}_i$  whenever  $d \in \mathcal{S}_i$ . Since the eigenvectors of  $H$  are identical to those of  $G$ , the subspaces  $\mathcal{S}_i$  are also invariant for  $H$ . Now consider  $S_i$  the orthogonal projectors onto  $\mathcal{S}_i$ . Since these projectors share a common system of eigenvectors with  $H$ , we know that they must commute, that is  $HS_i = S_iH$ . Using this very simple observation, we then obtain, for a step  $s_k$  at iteration  $k$  of a quasi-Newton algorithm, that

$$HS_i y_k = S_i H y_k = S_i s_k, \tag{10}$$

thereby yielding a new secant equation with the pair  $(S_i s_k, S_i y_k)$ . Repeating the procedure for  $i \in \{1, \dots, p\}$ , we therefore obtain  $p$  additional secant equations (in addition to the original equation (9)) provided we know the projections  $S_i$ .

If we now consider general twice differentiable, possibly non-convex, objective functions, then (9) remains valid for

$$G = \int_0^1 \nabla_{xx} f(x_k + t(x_{k+1} - x_k)) s_k dt$$

and the same reasoning still holds if the invariant subspaces are now associated with this latter matrix. Finally, if the subspaces  $\mathcal{S}_i$  are only approximately invariant, or if the operators  $S_i$  are only approximately equal to projectors onto these subspaces, then our secant equations stop being exact but can be expected

to hold approximately. We therefore refer to secant equations of the type (10) as *approximate*, as opposed to the *exact* equation (5). We are interested in the size of perturbations of  $G$  that would be necessary for the approximate secant equation (10) to hold exactly. We first consider all (possibly nonsymmetric) perturbations  $F_i$  such that

$$(G + F_i)S_i s_k = S_i y_k, \quad 1 \leq i \leq p. \quad (11)$$

A direct computation shows that for each  $i$ , the perturbation given by the Powell-Symmetric-Broyden (PSB) update (Powell [27])

$$E_i = \frac{(GS_i s_k - S_i y_k)(S_i s_k)^T + (S_i s_k)(GS_i s_k - S_i y_k)^T}{\|S_i s_k\|_2^2} - \frac{(GS_i s_k - S_i y_k)^T S_i s_k}{\|S_i s_k\|_2^2} \frac{S_i s_k S_k^T S_i^T}{\|S_i s_k\|_2^2}$$

satisfies (11) (with  $F_i = E_i$ ), and that  $\|E_i\|_2 \leq 3\|GS_i s_k - S_i y_k\|_2 / \|S_i s_k\|_2$ . Therefore,

$$\frac{\|E_i\|_2}{\|G\|_2} \leq 3 \frac{\|GS_i s_k - S_i y_k\|_2}{\|G\|_2 \|S_i s_k\|_2}, \quad (12)$$

for  $k \geq 0$  and  $0 \leq i \leq p$ . Interestingly, the factor 3 the right-hand side turns out to be unnecessary, as pointed out by Bunch, Demmel and Van Loan [4] (or Higham [18], p. 149), and we deduce that there exists a symmetric perturbation  $E_i$  such that  $E_i$  solves (11) and

$$\frac{\|E_i\|_2}{\|G\|_2} = \frac{\|GS_i s_k - S_i y_k\|_2}{\|G\|_2 \|S_i s_k\|_2}, \quad (13)$$

for  $k \geq 0$  and  $0 \leq i \leq p$ . If the quantity on the right-hand side of this inequality is modest, relation (10) ( $H = G^{-1}$ ) holds exactly for a symmetric Hessian approximation relatively close to  $G$ . It may thus be reasonable to use it for extracting curvature information, which may then be included in our current Hessian approximation. Let us now consider the numerator of the right-hand side of inequality (13), and assume that the operator  $S_i$  is symmetric and can be decomposed as  $S_i = Q_i D_i Q_i^T$ , where the columns of  $Q_i$  form an orthonormal basis of  $S_i$  and  $D_i$  is diagonal of dimension  $\dim(S_i)$ . Now consider  $Q_i$  such that the matrix  $(Q_i, Q_i^c)$  is orthogonal and define

$$Q_i^T G Q_i = G_i \quad \text{and} \quad (Q_i^c)^T G Q_i = F_i.$$

Then we have that

$$\begin{aligned} \|GS_i s_k - S_i y_k\|^2 &= \|G Q_i D_i Q_i^T s_k - Q_i D_i Q_i^T y_k\|^2 \\ &= \|G_i D_i Q_i^T s_k - D_i Q_i^T y_k\|^2 + \|F_i D_i Q_i^T s_k\|^2. \end{aligned}$$

Decomposing the step orthogonally as  $s_k = Q_i s_{ik} + Q_i^c s_{ik}^c$  and using the secant equation  $G s_k = y_k$ , we obtain that

$$\begin{aligned} \|GS_i s_k - S_i y_k\|^2 &= \|G_i D_i s_{ik} - D_i G_i s_{ik} - D_i F_i s_{ik}^c\|^2 + \|F_i D_i s_{ik}\|^2 \\ &\leq \|(G_i D_i - D_i G_i) s_{ik}\|^2 + \|F_i\|^2 \|D_i\|^2 (\|s_{ik}\|^2 + \|s_{ik}^c\|^2), \end{aligned}$$

and therefore, using (13) and the triangle inequality, that

$$\frac{\|E_i\|}{\|G\|} \leq \frac{\|G_i D_i - D_i G_i\|}{\sigma_{\min}(D_i) \|G\|} + \kappa(D_i) \frac{\|F_i\|}{\|G\|} \frac{\|s_k\|}{\|s_{ik}\|}, \quad (14)$$

where  $\sigma_{\min}(D_i)$  is the smallest singular value of  $D_i$ . This formula shows that we may expect the relative perturbation to  $G$  to be small when  $S_i$  is an approximate projector, in which case  $D_i$  is close to the identity matrix and of modest conditioning, and if the off-diagonal term  $F_i$  small compared to  $\|G\|$  (which we would expect if  $S_i$  is approximately invariant) together with  $\|s_{ik}\|$  being non-marginal with respect to  $\|s_k\|$ . This last condition is also acceptable, since we would not be interested in exploiting the curvature information along a “projected step”  $s_{ik}$  which is vanishingly small compared to the complete step  $s_k$ , because rounding errors would make then this information unreliable. Note that (14) also yields the further bound

$$\frac{\|E_i\|}{\|G\|} \leq \kappa(D_i) \left[ 2 \frac{\|G_i\|}{\|G\|} + \frac{\|F_i\|}{\|G\|} \frac{\|s_k\|}{\|s_{ik}\|} \right], \quad (15)$$

which indicates that, if  $\|G_i\|$  is small compared to  $\|G\|$ , we may also expect a relatively small perturbation to  $G$  even when  $D_i$  differs significantly from the identity (while remaining well-conditioned).

## 2.2. Multi-Secant (Limited-Memory) Quasi-Newton Algorithms

Using the simple derivation exposed above, and assuming, crucially, that a collection of (possibly approximate) projectors  $S_i$  are known, we may then outline a multi-secant quasi-Newton algorithm as in Algorithm 2.1.

### Algorithm 2.1: Multi-Secant Quasi-Newton Algorithm (outline)

- **Step 0: Initialization.** An initial point  $x_0 \in \mathbb{R}^n$  and an initial (positive-definite) Hessian  $H_0$  are given. The operators  $S_i$  ( $i = 1, \dots, p$ ) are also given, as well as a small tolerance  $\epsilon \geq 0$ . Compute  $f(x_0)$  and  $g(x_0)$ , and set  $k = 0$ .
- **Step 1:** If  $\|g(x_k)\| \leq \epsilon$ , stop.
- **Step 2: Compute the search direction.** Define  $d_k = -H_k g(x_k)$ .
- **Step 3: Linesearch.** Perform a linesearch ensuring

$$f(x_{k+1}) \leq f(x_k) + \alpha \langle g(x_k), d_k \rangle \quad \text{and} \quad \langle g(x_{k+1}), d_k \rangle \geq \beta \langle g(x_k), d_k \rangle$$

for some  $\alpha \in (0, 1)$  and  $\beta \in (\alpha, 1)$ , yielding  $f(x_{k+1})$ ,  $g(x_{k+1})$ ,  $s_k$  and  $y_k$  satisfying (8).

- **Step 4: Update the Hessian approximation.** Compute  $H_{k+1}$  such that (5) holds and

$$H_{k+1}(S_i y_k) = (S_i s_k) \quad i \in \{1, \dots, p\}. \quad (16)$$

- **Step 5: Loop.** Set  $k = k + 1$  and return to Step 1.

In this outline, we have not specified the linesearch procedure in detail, but this well understood technique is described in detail in Section 6.3 of Dennis and Schnabel [9] or in Moré and Thüente [23], for instance. We have not either indicated how we can impose (5) and (16) together. While updating simultaneously for more than one secant equation is indeed possible (see Byrd, Nocedal and Schnabel [7]), we focus here on the incorporation of the information in the Hessian in a sequential manner, by performing first  $p$  BFGS updates of the form (7) using the  $p$  pairs  $(S_i s_k, S_i y_k)$ , followed by a final update using the pair  $(s_k, y_k)$ . Of course the order in which the  $p$  “subspace updates” are performed is then significant.

If the problem size is large, which is the case of interest to us, it is often better to avoid the explicit updating of  $B_{k+1}$  or  $H_{k+1}$ , mostly because the low rank update (7) typically results in  $H_{k+1}$  being dense. To circumvent this problem, variants of the classical quasi-Newton algorithm have been proposed that replace (6) by a recursion in which  $H_{k+1}$  is implicitly reconstructed from some simple initial approximation (typically a multiple of the identity matrix: we use the matrix

$$\frac{\langle y_k, s_k \rangle}{\|y_k\|_2^2} I$$

in the experiments discussed below) and a modest number of the most recent pairs  $(s_j, y_j)$ . No matrix is ever assembled in the process, which only requires the storage of a small number of vectors. Because only the most recent pairs are used, these methods are called “limited-memory” methods, compared to the usual (full-memory) quasi-Newton algorithms where  $H_{k+1}$  includes information derived from all past pairs. The best-known method of this type is the limited-memory BFGS method pioneered by Byrd, Lu, Nocedal and Zhu [6]. The update (7) is (implicitly) used in this algorithm to compute the step. An efficient technique to perform the calculation of the step  $s_k$  is described on page 225 of Nocedal and Wright [26], where  $m$  pairs  $(s_j, y_j)$  are used sequentially in the implicit update of a diagonal initial approximation. We refer below to this technique as the implicit-secant-updating algorithm.

This technique can readily be adapted to our context where each secant pair generates  $p$  additional approximate ones corresponding to its images in  $p$  approximate invariant subspaces. Instead of considering only the last  $m$  pairs in the limited-memory updating procedure, we may now consider  $m \times p$  pairs, or any selection we care to make amongst them. An extreme case is when we select only the  $p + 1$  secant pairs corresponding to  $(s_k, y_k)$  and its images (16) onto  $\{S_i\}_{i=1}^p$ , giving a “memory-less” BFGS method.

### 2.3. Collinearity and Curvature Control

Nothing in the algorithm we have described so far prevents secant pairs from being linearly dependent. While not a major issue in the usual context where secant pairs are generated at different iterations, this might be an issue in our case, where one expects some dependency amongst the pairs generated from different, but possibly nested, invariant subspaces. Fortunately, some control of the possible collinearity of the secant pairs is easy, and we have chosen to include a provision in our algorithm that considers the angle between approximate secant pairs generated at iteration  $k$  and the “exact” pair  $(s_k, y_k)$ . More formally, we have decided to ignore the secant pair  $(S_i s_k, S_i y_k)$  whenever

$$|\langle S_i s_k, s_k \rangle|_2 \leq \tau \|S_i s_k\|_2 \|s_k\|_2 \quad \text{or} \quad |\langle S_i s_k, S_j s_k \rangle| \leq \tau \|S_i s_k\|_2 \|S_j s_k\|_2$$

for some  $\tau \in (0, 1]$  (typically 0.999 when this feature is active) and some  $j < i$ . Note that this last condition depends on the order in which the secant pairs are considered, a choice which we discuss below.

Because we wish to preserve the positive-definite nature of the approximate Hessian  $H_{k+1}$ , we also ignore approximate secant pairs for which

$$\langle S_i y_k, S_i s_k \rangle \geq \mu \|S_i y_k\|_2 \|S_i s_k\|_2$$

for some  $\mu \in (0, 1)$ .

Of course, the above discussion and the proposed multi-secant quasi-Newton methods may only be of practical interest when suitable operators  $S_i$  are known for a collection of (approximate) invariant subspaces. The purpose of the next paragraph is to show that this desirable situation may occur in at least an important practical case: multilevel optimization. But it should be noted that our derivation, and the bounds (14) and (15) are not restricted to this particular framework.

### 3. An Application to Multilevel Optimization

We now consider a particular framework in which the concepts described above occur naturally and therefore can be exploited to design improved optimization algorithms.

#### 3.1. Invariant Subspaces and Multigrid

Let us assume that the optimization variables of the problem under consideration represent the discretization of some continuous field defined on some spatial or temporal domain, a very common situation in engineering applications or physical modelling. For example, the variables may stand for coordinates of a design surface, atmospheric pressure over some part of the ocean, or position of a spacecraft along a controlled trajectory. The main characteristics of these problems is that it is possible to define discretization of the field of interest with varying degree of coarseness, from the very coarse to the very fine. For the sake of the argument, suppose that we consider  $r+1$  such different field discretizations, which we number from 0 (coarsest) to  $r$  (finest). In this case, it is very often reasonable to assume that, for each  $i \in \{1, \dots, r\}$ , there exist a full-rank linear operator  $R_i$  from  $\mathbb{R}^{n_i}$  into  $\mathbb{R}^{n_{i-1}}$  (the restriction from the fine grid  $i$  to the coarser grid  $i-1$ ) and another full-rank operator  $P_i$  from  $\mathbb{R}^{n_{i-1}}$  into  $\mathbb{R}^{n_i}$  (the prolongation from the coarse grid  $i-1$  to the fine grid  $i$ ) such that  $\sigma_i P_i = R_i^T$  for some known constant  $\sigma_i > 0$ . Moreover, these grid-transfer operators are typically computationally cheap to apply: the prolongation is for instance often chosen as the linear interpolation operator and the restriction as some multiple of its transpose, sometimes called the full-weighting operator.

When the problem to be solved for the (field) variables is a linear or nonlinear system of equations (instead of an optimization problem), multigrid techniques often yield the computationally most efficient algorithm, as their cost typically grows only linearly with the number of variables. The main characteristics of multigrid algorithms (we refer the reader to Briggs, Henson and McCormick [2] for an excellent introduction, or to Trottenberg, Oosterlee and Schüller [29], for a fuller treatment) are based on the observation that different “frequencies” are present in the solution of the finest grid problem (or even of the infinite-dimensional one), and become only progressively visible in the hierarchy from coarse to fine grids.



Low frequencies are visible from coarse grids and up, but higher ones can only be distinguished when the mesh-size of the grid becomes comparable to the inverse of the frequency in question. In multigrid strategies, specific algorithms, called *smoothers*, are known to very efficiently reduce the high frequency components of the error on a grid (that is, in most cases, the components whose “wavelength” is comparable to the grid’s mesh-size), but have little effect on the low frequency error components. It is observed however that such components on a fine grid appear more oscillatory on a coarser grid. They may thus be viewed as high frequency components on some coarser grid and be in turn efficiently reduced by a smoother. Moreover, this is done at a lower cost since computations on coarse grids are typically much cheaper than on finer ones. The multigrid strategy consists therefore in alternating between solving the problem on coarse grids, essentially annihilating low frequency components of the error, and on fine grids, where high frequency components are reduced (at a higher cost) by a smoother. This last operation is often called *smoothing* and the associated method a smoother because the effect of reducing high frequency components without altering much the low frequency ones has a “smoothing effect” on the error’s behaviour.

In other words, the multigrid strategy exploits the fact that the considered operator is (approximately) separable in the frequency domain, and that restrictions from fine to progressively coarser grids followed by prolongations to the fine grid isolate the corresponding nested invariant subspaces frequency-wise. A very well-understood example is that of the linear Poisson equation in a bounded domain, given in its discretized<sup>1</sup> variational form by

$$\min_u \frac{1}{2} \langle u, \Delta u \rangle - \langle f, u \rangle, \quad (17)$$

where  $u = u(x)$  is the unknown temperature distribution at position  $x$  of the underlying spatial domain and  $\Delta$  is the discretized Laplacian. It is easy (see Briggs, Henson and McCormick [2], page 18) to verify that the eigenvalues of the unidimensional discretized operator in  $n - 1$  variables are (when one takes the Dirichlet boundary condition  $u_0 = u_n = 0$  into account)

$$\lambda_i = 4 \sin^2 \left( \frac{i\pi}{2n} \right) \quad (i = 1, \dots, n - 1)$$

and that its eigenvectors are given componentwise by

$$z_{i,j} = \sin \left( \frac{ij\pi}{n} \right) \quad (i = 1, \dots, n - 1, j = 0, \dots, n).$$

It is remarkable that the eigenvectors  $z_i$  ( $i = 1, \dots, \frac{1}{2}n$ ) on the fine grid are exactly representable on a coarser grid with double mesh-size. If we choose the commonly used full weighting restriction operator (the transpose of the linear interpolator), it is possible to verify (see Briggs, Henson and McCormick [2], pages 80–81) that

$$[R_r z_i]_j = \theta_i \sin \left( \frac{ij\pi}{n/2} \right), \quad \text{where } \theta_i = \begin{cases} \cos^2 \left( \frac{i\pi}{2n} \right) & \text{for } i = 1, \dots, \frac{1}{2}n, \\ -\sin^2 \left( \frac{i\pi}{2n} \right) & \text{for } i = \frac{1}{2}n, \dots, n. \end{cases}$$

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<sup>1</sup>Using a simple finite-difference scheme.

It would be ideal if the prolongation  $P_r R_r z_i$  would be exactly a linear combination of  $\{z_i\}_{i=1}^{n/2}$ , the eigenvectors associated with the smooth modes on the fine grid, since then the image of  $P_r R_r$  would be identical to the invariant subspace spanned by the vectors  $\{z_i\}_{i=1}^{n/2}$ . Unfortunately, this is often not the case, due to a (often modest) contamination of the prolonged vector by oscillatory modes. Thus the image of  $P_r R_r$  is only approximately equal to  $\mathcal{S}_{r-1}$ , but this operator is typically very cheap to apply.

### 3.2. A Multi-Secant Limited-Memory BFGS Algorithm

We may now combine all the ingredients of our discussion so far into a single minimization algorithm: it suffices to use the multi-secant limited-memory BFGS method described above with the definition of our multigrid approximate (symmetric) operators

$$S_i u_r \stackrel{\text{def}}{=} P_r \dots P_{i+1} R_{i+1} \dots R_r u_r \quad \text{for } 0 \leq i \leq r-1. \quad (18)$$

In the usual multigrid framework, this operator would be combined with explicit smoothing steps in order to reduce the propagation of high frequency components in the result. More generally (and restricting the argument to two levels only and to the case where  $s_k$  belongs to the coarse subspace), the equation

$$R_i y_k = R_i G P_i P_i^+ s_k$$

also suggests to distinguish between differences in gradients (the vectors  $y_k$ ) to which the restriction operators  $R_i$  may naturally be applied (as gradients belong to the dual) and steps (the vectors  $s_k$ ) for which it would be more suitable to apply the generalized inverse  $P_i^+$  (these vectors lie in the primal). Unfortunately, these options are difficult to apply in practice because smoothing steps typically require the explicit knowledge of the Hessian matrix, which is unavailable here, and because computing the generalized inverse would be too costly. Moreover, as is discussed in Section 5, the limited-memory quasi-Newton method may itself be interpreted as a smoother, which thus provides an implicit *a posteriori* smoothing in the process.

We obviously expect multi-secant updates to work at their best potential when the eigensystem of the objective function's Hessian  $\nabla_{xx} f(x)$  is well-aligned with the grid, in the sense that

$$S_i \nabla_{xx} f(x) \approx \nabla_{xx} f(x) S_i.$$

Note that many algorithmic variants are possible in selecting  $\mathcal{P}_{k+1}$  (in Step 5). One may for instance give priority to the most recent information by selecting the pairs  $(s_{k,0}, y_{k,0}), \dots, (s_{k,r-1}, y_{k,r-1}), (s_k, y_k)$  or to exact secant equations (as in the usual limited-memory BFGS) by including pairs  $(s_{k-m+1}, y_{k-m+1}), \dots, (s_k, y_k)$  instead. Any combination of the above is also possible.

## 4. Numerical Experience with Multi-Secant Equations

We now illustrate the performance of the multi-secant multigrid limited-memory BFGS algorithm on test problems exhibiting the multigrid structure.

**Algorithm 3.1: Multi-Secant Multigrid Limited-Memory BFGS**

- **Initialization.** An initial point  $x_0 \in \mathbb{R}^{n_r}$  and an initial (positive-definite) Hessian  $H_0$  are given. The restriction and prolongation operators  $R_i$  and  $P_i$  are given for  $i = 1, \dots, r$ , and the operators  $S_i$  are given by (18) for  $i = 0, \dots, r-1$ . Choose a memory size  $m \geq 1$ , as well as a small tolerance  $\epsilon \geq 0$ . Compute  $f(x_0)$  and  $g(x_0)$ , define the initial set of secant pairs  $\mathcal{P}_0 = \emptyset$  and set  $k = 0$ .
- **Step 1:** If  $\|g(x_k)\| \leq \epsilon$ , stop.
- **Step 2: Compute the search direction.** Apply the implicit-secant-updating algorithm to compute  $s_k = -H_k g(x_k)$  using the secant pairs in  $\mathcal{P}_k$ .
- **Step 3: Linesearch.** Perform a linesearch ensuring

$$f(x_{k+1}) \leq f(x_k) + \alpha \langle g(x_k), s_k \rangle \quad \text{and} \quad \langle g(x_{k+1}), s_k \rangle \geq \beta \langle g(x_k), s_k \rangle$$

for some  $\alpha \in (0, 1)$  and  $\beta \in (\alpha, 1)$ , yielding  $f(x_{k+1})$ ,  $g(x_{k+1})$  and  $y_k$  satisfying (8).

- **Step 4: Generate secant pairs.** Apply the operators to compute

$$s_{k,i} = S_i s_k \quad \text{and} \quad y_{k,i} = S_i y_k$$

for  $i = 0, \dots, r-1$  and set

$$\mathcal{P}_k^+ = \mathcal{P}_k \cup \{(s_{k,0}, y_{k,0}), \dots, (s_{k,r-1}, y_{k,r-1}), (s_k, y_k)\}.$$

- **Step 5: Select the next set of secant pairs.** Select  $m$  pairs in  $\mathcal{P}_k^+$  to form  $\mathcal{P}_{k+1}$ .
- **Step 6: Loop.** Set  $k = k + 1$  and return to Step 1.

**4.1. Test Examples**

We now briefly describe the problems on which our multi-secant limited-memory BFGS algorithm has been applied.

*4.1.1. DN: A Dirichlet-to-Neumann Transfer Problem*

Let  $S$  be the square  $[0, \pi] \times [0, \pi]$  and let  $\Gamma$  be its lower edge defined by  $\{(x, y), 0 \leq x \leq \pi, y = 0\}$ . The Dirichlet-to-Neumann transfer problem (Lewis and Nash [21]) consists in finding the function  $a(x)$  defined on  $[0, \pi]$ , that minimizes

$$\int_0^\pi [\partial_y u(x, 0) - \phi(x)]^2 dx,$$

where  $\partial_y u$  is the partial derivative of  $u$  with respect to  $y$ , and where  $u$  is the solution of the boundary value problem

$$\begin{aligned} \Delta u &= 0 && \text{in } S, \\ u(x, y) &= a(x) && \text{on } \Gamma, \\ u(x, y) &= 0 && \text{on } \partial S \setminus \Gamma. \end{aligned}$$

The problem is a one-dimensional minimization problem, but the computations of the objective function and gradient involve a partial differential equation in

two dimensions. To introduce oscillatory components in the solution, we set  $\phi(x) = \sum_{i=1}^{15} [\sin(ix) + \sin(40x)]$ . The discretization of the problem is performed by finite differences with the same grid spacing in the two directions. The discretized problem is a linear least-squares problem.

#### 4.1.2. Q2D: A Simple Quadratic Example

We consider here the two-dimensional model problem for multigrid solvers in the unit square domain  $S_2 = [0, 1] \times [0, 1] = \{(x, y), 0 \leq x \leq 1, 0 \leq y \leq 1\}$ :

$$\begin{aligned} -\Delta u(x, y) &= f \text{ in } S_2 \\ u(x, y) &= 0 \text{ on } \partial S_2, \end{aligned}$$

where  $f$  is such that the analytical solution to this problem is  $u(x, y) = 2y(1 - y) + 2x(1 - x)$ . This problem is discretized using a 5-point finite-difference scheme, giving linear systems  $A_i x = b_i$  at level  $i$  where each  $A_i$  is a symmetric positive-definite matrix. The optimization is carried out on the variational minimization problem

$$\min_{x \in \mathbb{R}^{n_r}} \frac{1}{2} x^T A_r x - x^T b_r, \quad (19)$$

which is obviously equivalent to the linear system  $A_r x = b_r$ . The main purpose of this example is to illustrate that our algorithm is able to exploit the best multigrid structure.

#### 4.1.3. MS: A Minimum Surface Problem

We consider the minimum surface problem

$$\min_{v \in K} \int_0^1 \int_0^1 \sqrt{1 + (\partial_x v)^2 + (\partial_y v)^2} dx dy,$$

where  $K = \{v \in H^1(S_2) \mid v(x, y) = v_0(x, y) \text{ on } \partial S_2\}$ . The boundary condition  $v_0$  is chosen as

$$v_0(x, y) = \begin{cases} f(x), & y = 0, 0 \leq x \leq 1, \\ 0, & x = 0, 0 \leq y \leq 1, \\ f(x), & y = 1, 0 \leq x \leq 1, \\ 0, & x = 1, 0 \leq y \leq 1, \end{cases}$$

where  $f(x) = x(1 - x)$ . This convex problem is discretized using a finite element basis defined using a uniform triangulation of  $S_2$ , with same grid spacing  $h$  along the two coordinate directions. The basis functions are the classical P1 functions which are linear on each triangle and take value 0 or 1 at each vertex. The starting point is a random vector uniformly distributed in  $[0, 1]$ .

#### 4.1.4. BR: Bratu's Problem in Two Dimensions

We consider the minimization problem

$$\min_{v \in K} \int_0^1 \int_0^1 ((\partial_x v)^2 + (\partial_y v)^2 + R v e^v) dx dy,$$

where  $K = \{v \in H^1(S_2) \mid v(x, y) = 0 \text{ on } \partial S_2\}$ , corresponding to a variational formulation of the partial differential equation

$$\begin{aligned} -\Delta u(x, y) + Re^u &= 0 \text{ in } S_2 \\ u(x, y) &= 0 \text{ on } \partial S_2, \end{aligned}$$

where  $R = 6.8$ , as advocated in Moré and Toraldo [24]. This variational problem is discretized using the same finite element basis as that used in the MS problem. The starting point is a random vector uniformly distributed in  $[0, 1]$ .

#### 4.1.5. IP: An Inverse Problem from Image Processing

We consider the image deblurring problem stated on page 130 of Vogel [30]. In this problem, the columns of the unknown deblurred image are stacked into a vector  $f$ . A doubly block Toeplitz matrix  $T$  is computed using the `blur` function of Hansen's toolbox [17], which also yields the blurred image  $d$ . The image deblurring problem uses the total variation principle and can be written as

$$\min_f \left[ \frac{1}{2} \|Tf - d\|_2^2 + \frac{1}{10} \int_0^1 \int_0^1 \sqrt{1 + (\partial_x f)^2 + (\partial_y f)^2} dx dy \right].$$

The problem is convex. The discretization scheme is the same as for the MS problem, and the starting point for the minimization is chosen as  $f = 0$ .

#### 4.1.6. MT: A Membrane Tracking Problem

We consider the minimization problem

$$\min_{v \in K} \int_0^1 \int_0^1 (\partial_x v)^2 + (\partial_y v)^2 + v dx dy.$$

The set of constraints is defined by

$$K = \{v \in H^1(S_2) \mid v(x, y) = 0 \text{ on } \Gamma_1 \text{ and } v(x, y) \geq l(x, y) \text{ on } \Gamma_2\},$$

where boundary conditions are defined on the line segments  $\Gamma_1 = \{0\} \times [0, 1]$  and  $\Gamma_2 = \{1\} \times [0, 1]$ , and where  $l(x, y) = \sqrt{1 - (y - 0.5)^2} - 1.3$ . This problem corresponds to the displacement of a membrane under a traction of unit density (see Hlavacek, Haslinger, Necas and Lovisek [19]). To fit into our unconstrained optimization framework, this bound-constrained problem is transformed into a problem with Dirichlet boundary conditions by replacing  $K$  with

$$K' = \{v \in H^1(S_2) \mid v(x, y) = 0 \text{ on } \Gamma_1 \text{ and } v(x, y) = l(x, y) \text{ on } \Gamma_2\}.$$

The discretization of the above problem again uses the same finite element basis as for the MS problem. The starting point is a random vector uniformly distributed in  $[0, 1]$ .

#### 4.1.7. DASW: Data Assimilation for the Shallow-Water System

Data assimilation problems constitute an important class of parameter estimation. Their purpose is to reconstruct the initial conditions at  $t = 0$  of a dynamical system based on knowledge of the system evolution laws and on observations of the state at times  $t_i$ . More precisely, consider a dynamical system described by the equation  $\dot{x} = f(t, x)$  whose solution operator is given by  $x(t) = \mathcal{M}(t, x_0)$ . Assume

that the system state is observed (possibly only in parts) at times  $\{t_i\}_{i=0}^N$ , yielding observation vectors  $\{y_i\}_{i=0}^N$ , whose model is given by  $y_i = \mathcal{H}x(t_i) + \epsilon$ , where  $\epsilon$  is a noise with covariance matrix  $R_i$ . Assume finally that one knows  $B$ , an *a priori* error covariance matrix on  $x_0$ . We are then interested to find  $x_0$  which minimizes

$$\frac{1}{2}\|x_0 - x_b\|_{B^{-1}}^2 + \frac{1}{2}\sum_{i=0}^N \|\mathcal{H}\mathcal{M}(t_i, x_0) - y_i\|_{R_i^{-1}}^2. \quad (20)$$

The first term in this cost function is the often called the background term, the second the observation term.

An interesting application of (20) is given by a system governed by the (two-dimensional) shallow-water equations. This system is often considered as a good approximation of the dynamical systems used in ocean modeling, themselves a crucial element of climatic evolution scenarios (Griffies [16]). The system's equations are

$$\begin{cases} \frac{\partial u}{\partial t} + u\frac{\partial u}{\partial x} + v\frac{\partial v}{\partial y} - fv + g\frac{\partial z}{\partial x} &= 0, \\ \frac{\partial v}{\partial t} + u\frac{\partial v}{\partial x} + v\frac{\partial v}{\partial y} + fu + g\frac{\partial z}{\partial y} &= 0, \\ \frac{\partial z}{\partial t} + u\frac{\partial z}{\partial x} + v\frac{\partial z}{\partial y} + z\left(\frac{\partial u}{\partial x} + \frac{\partial v}{\partial y}\right) &= 0, \end{cases} \quad (21)$$

where  $u$ ,  $v$  and  $z$  are functions of  $(x, y, t)$ . The domain is the rectangle  $[0, L_x] \times [0, L_y]$  (with  $L_x = 32 \times 10^6$  meters and  $L_y = 8 \times 10^6$  meters) and the integration horizon is 50 timesteps of 400 seconds. The boundary conditions are assumed to be periodic in  $y$  and of Dirichlet type in  $x$ . Following the suggestion by Weaver and Courtier [31] we have modelled the *a priori* term  $x_b$  using a diffusion operator. As is recommended in the climate modelling community (we refer the reader to Griffies [16], for further details), our formulation uses initial geostrophic winds and a  $\beta$ -plane formulation for the Coriolis force; we integrate this system using a leapfrog scheme and a Laplacian spatial damping (which introduces a right-hand side of the form  $\mu(\Delta u, \Delta v, \Delta z)^T$  in the system (21), where  $\mu = 3 \times 10^6$ ). We also considered using an Asselin time-filter (see Asselin [1]), but gave this up as it had little effect on our results. In our problem, we assume to observe the state at every fifth point in the spatial domain and every fifth time step. The true initial geopotential height, which we seek to reconstruct, is assumed to be

$$z(x, y, 0) = 5000 + 50 \sin\left(\frac{2\pi x}{L_x}\right) \cos^2\left(\frac{\pi(y - \frac{1}{2}L_y)}{L_y}\right).$$

The same starting point for this non-convex problem is used for all tests and is generated as the initial state perturbed by normally distributed random noise with zero mean and standard deviation equal to  $10^{-2}$ .

This problem is interesting for our tests because numerical experiments indicate that this problem is not well-suited to a multigrid approach<sup>1</sup>.

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<sup>1</sup>We experimented with a Galerkin full-multigrid scheme on a linear system whose matrix was obtained from finite differences in the gradients of (20), and observed a numerical performance not significantly better than that of a pure Gauss-Seidel smoother.

#### 4.2. The First Results

We now turn to the numerical experiments themselves, where we considered the test examples described above for the sizes and algorithmic parameters presented in Table 1. In this table,  $m$  is the total number of secant pairs (exact and approximate) memorized by our algorithm and “# levels” is the number of levels exploited in the multi-secant algorithm. All experiments were run in Matlab. We used the full-weighting restriction and the linear-interpolation prolongation operators in all cases, and  $\mu$  was set to  $10^{-6}$ . Convergence of the algorithm was declared as soon as  $\|g(x_k)\| \leq \epsilon$ , the values of this tolerance being also specified in Table 1.

		DN	Q2D	MS	BR
size	( $n$ )	255	$1046529 = 1023^2$	$1046529 = 1023^2$	$261121 = 511^2$
memory	( $m$ )	10	9	9	7
# levels	( $r$ )	7	8	7	3
accuracy	( $\epsilon$ )	$10^{-5}$	$10^{-5}$	$10^{-5}$	$10^{-5}$

		IP	MT	DASW
size	( $n$ )	$66049 = 257^2$	$262143 = 511 \times 513$	$3969 = 63^2$
memory	( $m$ )	9	8	5
# levels	( $r$ )	4	6	4
accuracy	( $\epsilon$ )	$10^{-5}$	$10^{-5}$	$10^{-3}$

Table 1. The parameters in our numerical tests

The algorithmic variants tested differ according to the mechanism used in Step 5 to select the secant pairs, the order in which these pairs are used in the implicit updating procedure and the level of collinearity control  $\tau$ . We have defined three different strategies for the selection of secant pairs:

- **full:** all secant pairs (exact and approximate) are considered for each iteration, but the pairs generated at the iterations further in the past are dropped first, and, amongst those corresponding to the same iteration, the approximate pairs are dropped before the exact one;
- **local:** only the approximate secant pairs generated at the current iterations are considered for updating, in addition to as many past exact pairs as allowed by the memory;
- **mless:** all information from previous iterations is discarded and only the (exact and approximate) pairs generated at the current iteration are considered. This corresponds to using the multiple-secant updates in a purely “memoryless” manner.

We have also defined two different strategies for the order in which the approximate secant pairs are used for updating :

- **coarse-first:** the approximate inverse Hessian is updated for the pairs corresponding to the coarser levels first;
- **fine-first:** the approximate inverse Hessian is updated for the pairs corresponding to the finer levels first.

Each algorithmic variant is thus characterized by the triplet specifying its the pair selection strategy, its pair ordering strategy and its value of  $\tau$ . We also consider the standard limited-memory BFGS method (L-BFGS) for comparison. Table 2 on the next page presents the results of our first tests, expressed in terms of number of objective function evaluations (nf) and iterations (nit).

These results indicate that using approximate secant pairs associated with grid levels is potentially useful, although not uniformly for every problem nor across

Multi-secant LM variant			DN		Q2D		MS		BR	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	152	122	381	304	2130	1716	472	381
full	coarse-first	0.999	131	108	515	432	2541	2194	547	465
full	fine-first	1.0	130	105	366	273	2097	1815	428	367
full	fine-first	0.999	120	98	638	550	2361	2099	396	337
local	coarse-first	1.0	94	84	563	427	2003	1707	440	372
local	coarse-first	0.999	110	92	501	385	2033	1725	450	384
local	fine-first	1.0	120	100	304	233	2267	1943	397	339
local	fine-first	0.999	90	76	335	278	1867	1605	400	341
mless	coarse-first	1.0	125	100	563	427	1724	1364	527	422
mless	coarse-first	0.999	113	89	501	385	2207	1832	401	316
mless	fine-first	1.0	137	100	304	233	1679	1371	420	313
mless	fine-first	0.999	140	107	335	278	2204	1844	445	338
L-BFGS			330	319	1505	1471	2671	2644	1074	1053

Multi-secant LM variant			IP		MT		DASW	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	196	186	1542	1277	73	70
full	coarse-first	0.999	204	206	1262	1041	73	70
full	fine-first	1.0	213	200	1070	883	64	62
full	fine-first	0.999	214	210	1272	1077	64	62
local	coarse-first	1.0	165	161	929	781	73	70
local	coarse-first	0.999	165	162	1070	882	73	70
local	fine-first	1.0	185	184	1027	861	64	62
local	fine-first	0.999	180	179	1138	959	64	62
mless	coarse-first	1.0	355	353	1536	1221	73	70
mless	coarse-first	0.999	280	279	1307	1051	73	70
mless	fine-first	1.0	265	256	862	650	64	62
mless	fine-first	0.999	232	225	1461	1149	64	62
L-BFGS			181	176	1419	1386	60	57

Table 2. Performance of multi-secant limited-memory BFGS algorithms

variants. Further variations were also observed for other choices of  $m$  and  $r$ , but leave the overall picture unchanged. In the results reported here, the improvement is especially noticeable for problems (like Q2D and BR) where one expects the multigrid method to perform well. One also notes that the “local” and “mless” variants seem to be most efficient, and that collinearity control often helps somewhat the first of these strategies. The memoryless variants (“mless”) are especially efficient on the Laplacian (Q2D) and Bratu (BR) problems. We also note that the memory-less and local variants give identical number of function evaluations and iterations for problem Q2D, which we believe results from the fact that  $r \approx m$ , which leaves little room for past iteration history. The same comment applies to the data assimilation problem (DASW), for which the multiseant variants are clearly less successful. Given our observation that this latter example is not well-suited to multigrids, the performance of our algorithms on this problem does not come as a surprise. We finally observe that collinearity control produces fairly mixed results, and seems to be irrelevant for problem DASW.

In order to verify the analysis of Section 2.1, we have also computed the relative perturbations (13) (with  $G = \nabla_{xx}f(x_{k+1})$  and  $\|\cdot\| = \|\cdot\|_\infty$ ) during runs of the variant [local,coarse-first,1.0] on problems Q2D and MS. Typical results are shown in Figure 1.

This figure shows that the size of the relative perturbation of the true Hessian needed to make the approximate secant exact is very modest (a few percent of  $\|\nabla_{xx}f(x_{k+1})\|$ , typically). The relative size of the Hessian perturbation necessary make the exact secant equation (5) hold exactly is shown in the right figure as the



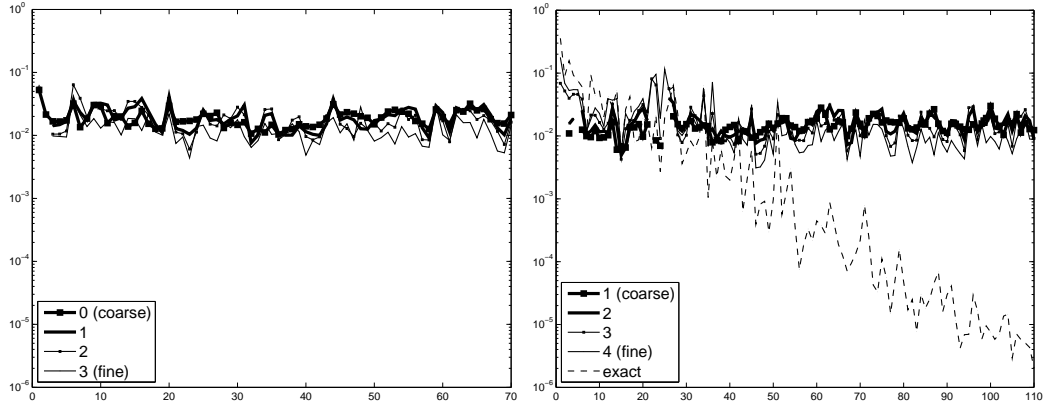


Figure 1. Evolution of the relative Hessian perturbation sizes (13) with  $k$  for the variant [local,coarse-first,1.0] with  $m = 7$  and  $r = 4$  for problems of size  $n = 63^2$  (Q2D on the left, MS on the right, logarithmic vertical scale)

curve ultimately decreasing to the order of  $10^{-5}$ . Clearly, this perturbation is of the same order as that for the approximate secant equations in the early iterations of the algorithm. It is invisible on the left figure, because it is always tiny (between  $10^{-15}$  and  $10^{-10}$ ) on a quadratic function. Further analysis (not illustrated here) indicates that the perturbation corresponding to past exact secant equations follows the same pattern as that corresponding to the current one, both for quadratic and nonquadratic problems.

## 5. Asymptotic Approximate Eigenvalue Equations

These numerical experiments prompt another discussion. It was observed in the numerical test-runs that the multiple-secant limited-memory BFGS algorithm also acts as a smoother on the original problem, in the sense that convergence often occurs much faster for the oscillatory modes of the solution than for the smooth modes. This is illustrated in Figure 2 which shows the decomposition of the step  $s_k$  along the subspaces  $\mathcal{S}_0, \mathcal{S}_1^C, \mathcal{S}_2^C, \dots, \mathcal{S}_r^C$ , where, for  $i > 0$ ,  $\mathcal{S}_i^C$  is the orthogonal complement of  $\mathcal{S}_{i-1}$  in  $\mathcal{S}_i$ , thereby isolating to contribution which is specific to each of the nested subspaces.

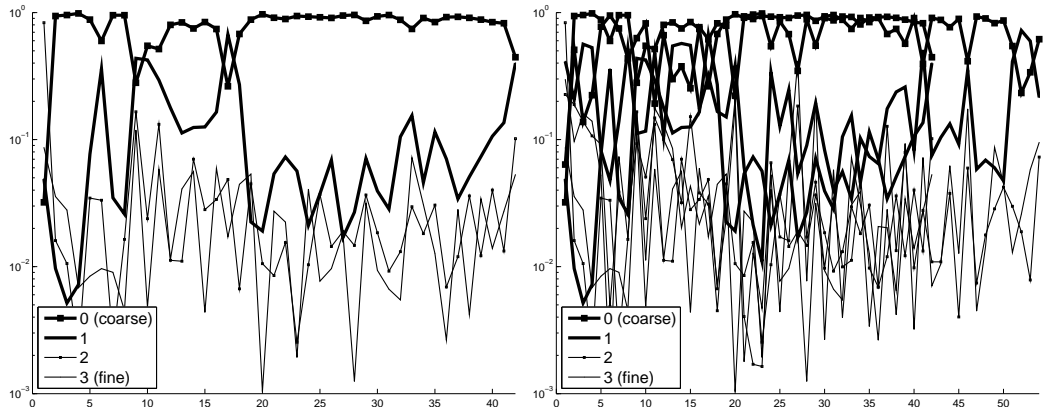


Figure 2. Evolution of the decomposition of the step  $s_k$  for the variant [local,coarse-first,1.0] with  $m = 7$  and  $r = 4$  along  $\mathcal{S}_0, \mathcal{S}_1^C, \mathcal{S}_2^C, \dots, \mathcal{S}_r^C$  for problems of size  $n = 31^2$  (Q2D on the left, MS on the right, logarithmic vertical scale)

In this figure, one observes that, for both problems,  $s_k$  is nearly entirely contained in  $\mathcal{S}_0$  (the coarse subspaces corresponding to very smooth modes), for  $k$  sufficiently

large, which we express by writing that

$$s_k \approx S_0 s_k \quad (22)$$

for large  $k$ . Now observe that the eigenvalues associated with  $\mathcal{S}_0$  are typically the smallest ones: in the case of the one-dimensional Laplacian model problem analyzed above, we have that

$$S_0 = \text{span} \left\{ \sin \left( \frac{ij\pi}{n} \right) \right\}_{i=1}^{n_0}$$

and the eigenvalues associated with this invariant subspace are

$$\lambda_i = 4 \sin^2 \left( \frac{i\pi}{2n} \right) \quad (i = 1, \dots, n_0). \quad (23)$$

If we assume, for instance, that  $n = 255$  and that we exploit five levels, we verify that  $n_0 = 7$ . The eigenvalues (23) are therefore all contained in the interval  $[0.00015, 0.0074]$ . The norm of the Laplacian operator is however given by its maximal eigenvalue, equal to 4. Thus the deviation of the eigenvalues associated to  $\mathcal{S}_0$  from any approximation  $\rho_0$  in the interval, relative to the operator norm, is bounded by

$$\frac{0.0074 - 0.00015}{4} \approx 0.0018,$$

that is slightly less than 0.2%. Modifying the Laplacian operator by imposing that its restriction to  $\mathcal{S}_0$  is  $\rho_0 I$  therefore amounts to a perturbation of the operator of at most 0.2%, which is again very modest. Moreover, this discussion is consistent with our observation following (15) as we see here that  $\|G_0\|$  is small compared to  $\|G\|$ . Returning to the case discussed in Section 2.1, this indicates that

$$Gs_k \approx GS_0 s_k \approx \rho_0 s_k,$$

and  $s_k \approx S_0 s_k$  is an approximate eigenvector of  $G$ , the value  $\rho_0$  being the Rayleigh quotient along this direction, given by

$$\rho_0 \|S_0 s_k\|^2 = \langle S_0 s_k, GS_0 s_k \rangle \approx \langle S_0 s_k, S_0 G s_k \rangle \approx \langle S_0 s_k, S_0 y_k \rangle.$$

Moreover,

$$S_0 y_k = S_0 G s_k = GS_0 s_k \approx G s_k = y_k. \quad (24)$$

If we now consider  $\mathcal{S}_1$  (a superset of  $\mathcal{S}_0$  of dimension 15 in our example), the size of relative perturbation Hessian remains at most 9%. Expressing the above reasoning for an a generic invariant subspace  $\mathcal{S}_i$  (note that  $\mathcal{S}_0 \subset \mathcal{S}_i$  in our multilevel context), we may thus expect that

$$GS_i s_k \approx \frac{\langle GS_i s_k, S_i s_k \rangle}{\|S_i s_k\|_2^2} S_i s_k \approx \frac{\langle y_k, S_i s_k \rangle}{\|S_i s_k\|_2^2} S_i s_k \approx \frac{\langle S_i y_k, S_i s_k \rangle}{\|S_i s_k\|_2^2} S_i s_k,$$

which yields

$$HS_i s_k \approx \frac{\|S_i s_k\|^2}{\langle S_i y_k, S_i s_k \rangle} S_i s_k. \quad (25)$$

Alternatively,

$$\langle S_i y_k, HS_i y_k \rangle \approx \langle S_i y_k, S_i H y_k \rangle \approx \langle S_i y_k, S_i s_k \rangle \approx \langle y_k, S_i s_k \rangle,$$

yielding

$$HS_i s_k \approx \frac{\langle S_i y_k, S_i s_k \rangle}{\|S_i y_k\|_2^2} S_i s_k. \quad (26)$$

Using the approximations (22) and (24), we see that equations of the form

$$Hp_{i,k} \approx \rho_{i,k} q_{i,k} \quad (i = 0, \dots, r-1), \quad (27)$$

may therefore be of interest when updating our Hessian approximation, where we have the choice to select  $p_{i,k}$  and  $q_{i,k}$  among  $S_i s_k$  ( $0 \leq i < r-1$ ) or  $s_k$ , and  $\rho_{i,k}$  among

$$\frac{\|S_i s_k\|_2^2}{\langle y_k, S_i s_k \rangle}, \quad \frac{\|S_i s_k\|_2^2}{\langle S_i y_k, S_i s_k \rangle}, \quad \frac{\langle y_k, S_i s_k \rangle}{\|S_i y_k\|_2^2} \quad \text{or} \quad \frac{\langle S_i y_k, S_i s_k \rangle}{\|S_i y_k\|_2^2} \quad (28)$$

Further alternative forms may be obtained by replacing  $S_i s_k$  by  $s_k$  or  $\|S_i s_k\|_2^2$  by  $\langle s_k, S_i s_k \rangle$  or  $\|s_k\|_2^2$ , or  $S_i y_k$  by  $y_k$  in (28)...

We also note that Algorithm 3.1 was stated above using only secant equations of the form (16), but it is now possible to incorporate equations of the form (27) (possibly in combination with (16)) into the updates. This only requires the redefinition of the vectors  $y_{k,i}$  to be generated at Step 4: for instance, we redefine  $y_{k,i}$  to be

$$\frac{\langle y_{k,i}, s_{k,i} \rangle}{\|s_{k,i}\|_2^2} s_{k,i}$$

if one wishes to use (25).

Finally, we observe that the variants described in this section differ from those in Section 2.1 in that they only require the storage of the steps  $s_k$ . This significantly reduces the memory needs for the method, which might be an advantage for very large problems.

## 6. Numerical Experience with Eigenvalue Equations

Our present purpose is not to conduct a full investigation of which of the alternative forms mentioned in the previous section is numerically preferable and when. We focus in this section on verifying the main lines of our analysis of approximate eigenvalue equations and on illustrating the potential of a few choices in a simple algorithmic setting. We leave the more complete experimentation and the design of a suitable selection mechanism for further research.

In this first exploration, we have chosen to experiment with four variants of the many possible choices of the type (27)-(28), namely the ‘‘eigenvalue equations’’

given by (25), (26),

$$HS_i s_k \approx \frac{\|S_i s_k\|_2^2}{\langle y_k, S_i s_k \rangle} S_i s_k, \quad (29)$$

and

$$HS_i s_k \approx \frac{\langle y_k, S_i s_k \rangle}{\|S_i y_k\|_2^2} S_i s_k. \quad (30)$$

The results of the numerical experience using these equations in the limited memory framework already detailed for approximate secant equations are given in Tables 3 to 9.

Eigenvalue LM variant			(29)		(25)		(30)		(26)	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	219	129	158	126	132	118	139	114
full	coarse-first	0.999	222	125	129	106	177	160	153	129
full	fine-first	1.0	529	272	176	134	129	117	145	106
full	fine-first	0.999	243	127	157	118	145	129	113	89
local	coarse-first	1.0	156	96	121	101	233	217	138	117
local	coarse-first	0.999	164	106	122	99	195	178	126	103
local	fine-first	1.0	159	100	104	86	116	104	101	85
local	fine-first	0.999	159	103	127	105	135	124	123	105
mless	coarse-first	1.0	199	112	169	136	246	223	138	111
mless	coarse-first	0.999	224	128	178	132	239	219	153	125
mless	fine-first	1.0	229	128	122	90	229	192	168	135
mless	fine-first	0.999	216	119	155	117	188	159	156	125

Table 3. Performance of eigenvalue limited-memory BFGS algorithms on problem DN (L-BFGS uses 330 function evaluations and 319 iterations on this problem)

Eigenvalue LM variant			(29)		(25)		(30)		(26)	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	707	564	806	608	689	534	552	440
full	coarse-first	0.999	592	470	868	672	704	592	689	595
full	fine-first	1.0	414	316	439	336	565	467	842	675
full	fine-first	0.999	908	749	1068	853	1185	966	770	637
local	coarse-first	1.0	551	437	771	576	630	482	686	551
local	coarse-first	0.999	405	313	524	410	532	410	602	481
local	fine-first	1.0	690	528	479	366	558	433	520	428
local	fine-first	0.999	579	479	389	313	699	582	518	413
mless	coarse-first	1.0	551	437	771	576	630	482	686	551
mless	coarse-first	0.999	405	313	524	410	532	410	532	481
mless	fine-first	1.0	690	528	479	366	558	433	520	428
mless	fine-first	0.999	579	479	389	313	699	582	518	413

Table 4. Performance of eigenvalue limited-memory BFGS algorithms on problem Q2D (L-BFGS uses 1505 function evaluations and 1471 iterations on this problem)

Qualitatively speaking, we obtain a conclusion very similar to that reached for approximate secant equations: there exists a clear algorithmic potential in using the information specified by the approximate eigenvalue equations, but this potential is uniform neither across problems nor across algorithmic variants. We only note that some of the best performances are obtained with these techniques, in particular for the data assimilation problem DASW.

As for the multiple secant equations, we may complete our picture by an estimation of the approximate nature of the considered eigenvalue equations. This

Eigenvalue LM variant			(29)		(25)		(30)		(26)	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	2123	1672	2663	2027	2277	1998	2627	2326
full	coarse-first	0.999	2674	2145	2780	2170	2595	2295	2466	2187
full	fine-first	1.0	2300	1895	2192	1772	1957	1683	2101	1817
full	fine-first	0.999	3097	2561	2610	2117	3042	2588	3259	2792
local	coarse-first	1.0	2780	2270	2869	2227	2511	2213	2375	2135
local	coarse-first	0.999	2565	2094	2275	2214	2369	2078	2367	2099
local	fine-first	1.0	2147	1872	2168	1851	2147	1917	2029	1808
local	fine-first	0.999	1698	1470	1637	1379	1848	1654	2316	2080
mless	coarse-first	1.0	2840	2107	2246	2246	2356	2043	2180	1931
mless	coarse-first	0.999	2429	1889	2753	1997	2346	1997	2546	2182
mless	fine-first	1.0	1494	1207	1833	1477	1992	1718	1853	1643
mless	fine-first	0.999	2036	1974	1922	1627	3075	2674	3397	3045

Table 5. Performance of eigenvalue limited-memory BFGS algorithms on problem MS (L-BFGS uses 2671 function evaluations and 2644 iterations on this problem)

Eigenvalue LM variant			(29)		(25)		(30)		(26)	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	511	425	532	443	555	484	552	490
full	coarse-first	0.999	424	352	580	478	471	405	622	544
full	fine-first	1.0	321	251	420	326	469	392	593	494
full	fine-first	0.999	377	291	557	460	413	342	573	479
local	coarse-first	1.0	383	334	584	486	494	441	390	348
local	coarse-first	0.999	427	358	647	522	523	462	363	322
local	fine-first	1.0	365	329	548	466	527	465	315	282
local	fine-first	0.999	381	339	449	396	419	364	415	359
mless	coarse-first	1.0	660	525	518	391	575	473	548	468
mless	coarse-first	0.999	405	314	432	320	605	486	490	404
mless	fine-first	1.0	432	343	535	407	419	339	547	431
mless	fine-first	0.999	301	234	428	330	338	282	665	522

Table 6. Performance of eigenvalue limited-memory BFGS algorithms on problem BR (L-BFGS uses 1074 function evaluations and 1053 iterations on this problem)

Eigenvalue LM variant			(29)		(25)		(30)		(26)	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	330	303	282	274	256	252	248	241
full	coarse-first	0.999	367	314	206	199	237	229	233	228
full	fine-first	1.0	2007	1015	186	181	610	609	189	187
full	fine-first	0.999	1738	881	178	609	610	609	198	196
local	coarse-first	1.0	223	220	164	174	166	165	165	164
local	coarse-first	0.999	229	226	164	163	169	168	169	168
local	fine-first	1.0	680	405	179	178	211	210	184	183
local	fine-first	0.999	727	444	184	183	219	218	184	183
mless	coarse-first	1.0	285	261	195	189	344	343	191	189
mless	coarse-first	0.999	371	314	246	195	344	343	200	197
mless	fine-first	1.0	959	484	198	245	3061	3060	236	235
mless	fine-first	0.999	1292	695	240	239	3061	3060	233	230

Table 7. Performance of eigenvalue limited-memory BFGS algorithms on problem IP (L-BFGS uses 181 function evaluations and 176 iterations on this problem)

estimation is again obtained by measuring the size of the Hessian perturbation that would ensure (27) exactly, relative to the Hessian norm. In this case, the relative perturbation size is then given by

$$\frac{\|E_i\|_\infty}{\|G\|_\infty} \leq \frac{\|Gp_{i,k} - \rho_{i,k}q_{i,k}\|_\infty}{\|p_{i,k}\|_\infty \|G\|_\infty} \quad (31)$$

Eigenvalue LM variant			(29)		(25)		(30)		(26)	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	1219	1060	1363	1099	1219	1072	1292	1140
full	coarse-first	0.999	1370	1167	1295	1034	1588	1384	1469	1294
full	fine-first	1.0	974	805	935	749	1063	934	1232	1075
full	fine-first	0.999	1064	847	1334	1074	1605	1377	1593	1367
local	coarse-first	1.0	1544	1299	1221	987	1334	1190	1340	1205
local	coarse-first	0.999	1512	1279	1275	1033	1279	1124	1277	1128
local	fine-first	1.0	1129	964	977	823	1293	1113	1205	1043
local	fine-first	0.999	1262	1079	931	786	1312	1146	1496	1375
mless	coarse-first	1.0	1311	1050	1515	1531	1311	1300	1469	1256
mless	coarse-first	0.999	1304	1050	1682	1292	1775	1471	1417	1188
mless	fine-first	1.0	1106	881	1145	880	1423	1136	1064	866
mless	fine-first	0.999	890	703	1349	1094	1981	1649	1540	1270

Table 8. Performance of eigenvalue limited-memory BFGS algorithms on problem MT (L-BFGS uses 1419 function evaluations and 1386 iterations on this problem)

Eigenvalue LM variant			(29)		(25)		(30)		(26)	
Pairs	upd. order	$\tau$	nf	nit	nf	nit	nf	nit	nf	nit
full	coarse-first	1.0	65	62	83	80	47	45	79	71
full	coarse-first	0.999	65	62	83	80	47	45	79	71
full	fine-first	1.0	61	58	93	92	81	76	59	51
full	fine-first	0.999	61	58	93	92	81	76	59	51
local	coarse-first	1.0	65	63	83	80	47	45	79	71
local	coarse-first	0.999	65	63	83	80	47	45	79	71
local	fine-first	1.0	61	58	91	88	91	84	57	51
local	fine-first	0.999	61	58	91	88	91	84	57	51
mless	coarse-first	1.0	65	63	83	80	47	45	79	71
mless	coarse-first	0.999	65	63	83	80	47	45	79	71
mless	fine-first	1.0	61	58	91	88	91	84	57	51
mless	fine-first	0.999	61	48	91	88	91	84	57	51

Table 9. Performance of eigenvalue limited-memory BFGS algorithms on problem DASW (L-BFGS uses 60 function evaluations and 57 iterations on this problem)

for the appropriate choices of  $p_{i,k}$ ,  $q_{i,k}$  and  $\rho_{i,k}$ . The results obtained during runs of the variant [local,coarse-first,1.0] on problems Q2D and MS are illustrated in Figures 3 to 6. These pictures confirm our analysis that the limited-memory framework may generate useful approximate eigenvalue information. The quality and relevance of this information varies according to the precise choice of approximate eigenvalue equation used and problem considered. Once more, this encourages further investigation in methods exploiting this information .

## 7. Conclusions

We have shown that the *a priori* knowledge of approximate invariant subspaces associated with the Hessian of an unconstrained optimization problem allows a more efficient exploitation of the secant information, and thus more efficient minimization algorithms. As an application, we have also described how this knowledge can often be extracted from the multilevel structure of discretized infinite dimensional problems. Using the asymptotic smoothing properties of the limited-memory BFGS method, we have also indicated how the steps generated by this algorithm often generate approximate eigenvalue information, which may in turn be used to reduce the memory required for an efficient optimization. Preliminary numerical experience on a small collection of test problems suggests that this approach is promising and further investigation worthwhile.

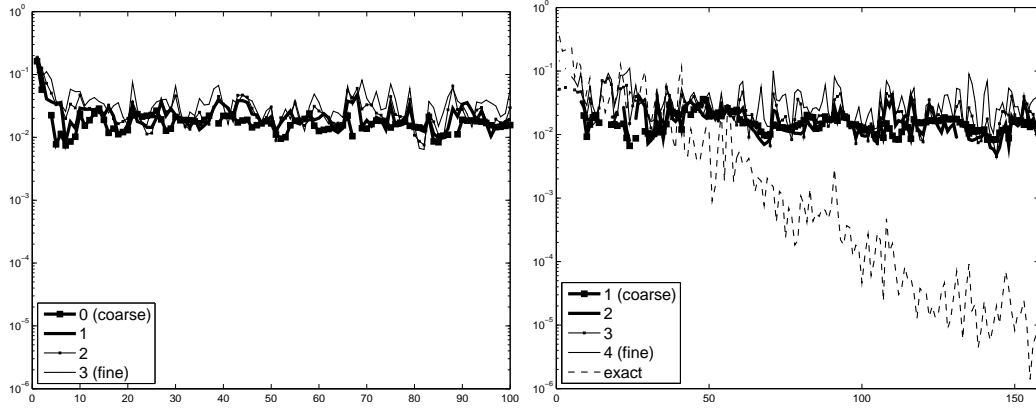


Figure 3. Evolution of the relative Hessian perturbation sizes (13) with  $k$  for the variant [local,coarse-first,1.0] with  $m = 7$  and  $r = 4$  using equation (29) for problems of size  $n = 63^2$  (Q2D on the left, MS on the right, logarithmic vertical scale)

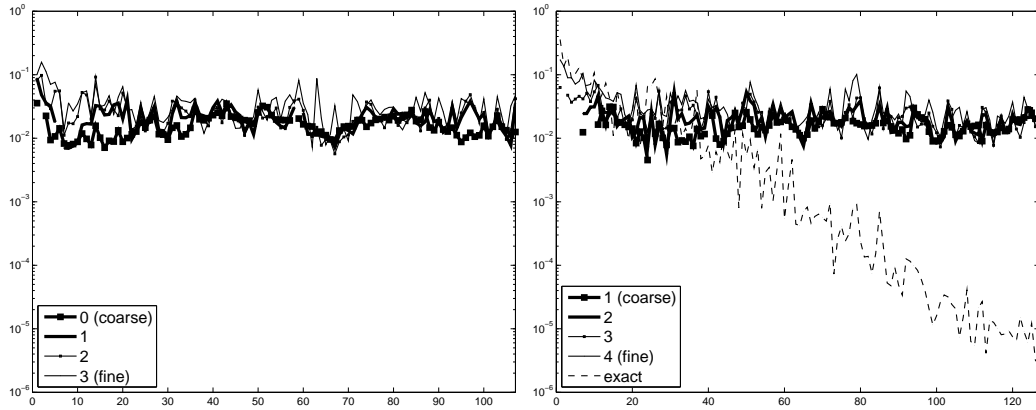


Figure 4. Evolution of the relative Hessian perturbation sizes (13) with  $k$  for the variant [local,coarse-first,1.0] with  $m = 7$  and  $r = 4$  using equation (25) for problems of size  $n = 63^2$  (Q2D on the left, MS on the right, logarithmic vertical scale)

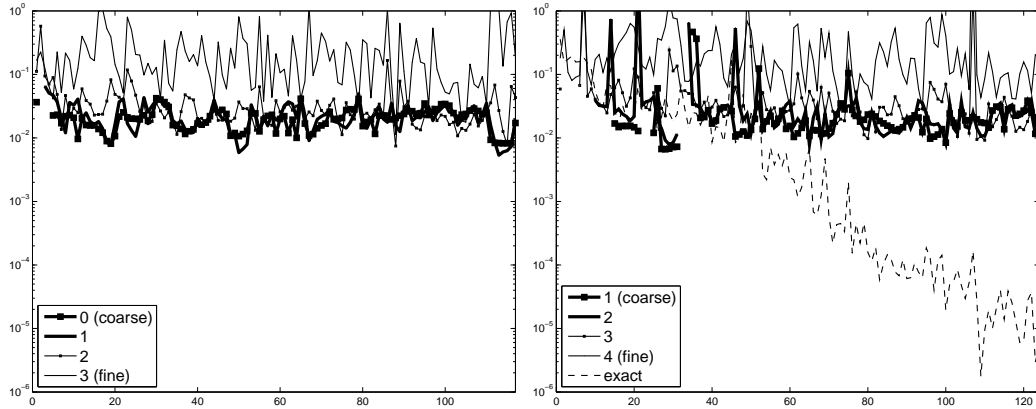


Figure 5. Evolution of the relative Hessian perturbation sizes (13) with  $k$  for the variant [local,coarse-first,1.0] with  $m = 7$  and  $r = 4$  using equation (30) for problems of size  $n = 63^2$  (Q2D on the left, MS on the right, logarithmic vertical scale)

The exploitation of invariant subspace information opens, from the authors' point of view, a number of immediate possibilities and more long-term perspectives. The first would be to investigate other frameworks where approximate invariant subspaces and associated approximate projectors can be obtained, possibly at the cost of some problem preprocessing. In particular the domain of model-reduction methods (of which multilevel can be considered a member) seems of interest. In

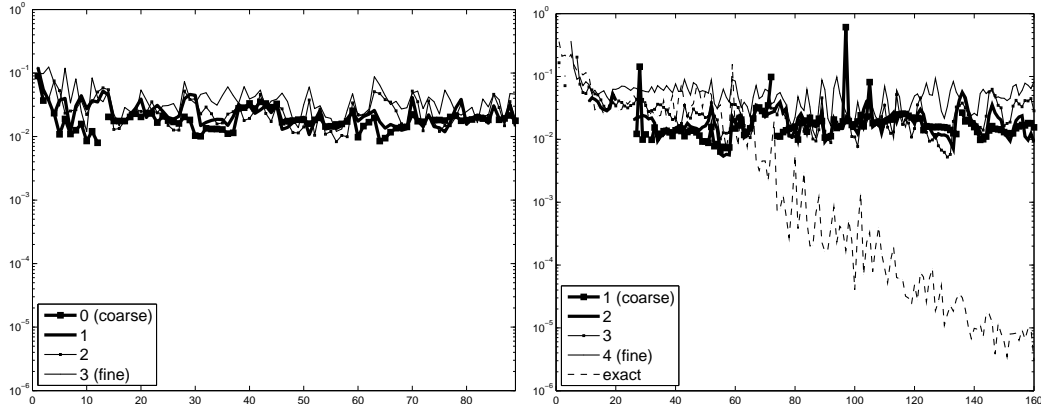


Figure 6. Evolution of the relative Hessian perturbation sizes (13) with  $k$  for the variant [local,coarse-first,1.0] with  $m = 7$  and  $r = 4$  using equation (26) for problems of size  $n = 63^2$  (Q2D on the left, MS on the right, logarithmic vertical scale)

this context, the hierarchy of models of different fidelity provides the support of the approximate invariant subspaces, and the authors believe that the ideas presented in this paper can then be applied using suitable prolongations and restrictions between the levels of this hierarchy. More algebraic decompositions may also be considered for the determination of approximate invariant subspaces, for instance in cases where nested subsets of variables are known to generate different degrees of model nonlinearity. One may also consider the use of multi-secant/eigenvalue approaches in algorithms for bound- and more generally constrained optimization.

In the more specific case of multilevel optimization and, although algorithmic variants have been outlined in the present paper, much remains to be done for obtaining a well-tested, robust and optimized multi-secant/eigenvalue quasi-Newton code. In particular, approximate techniques for computing the vector  $P_i^+ s_k$  as an alternative to  $R_i s_k$  might prove efficient in this context. But other questions also merit further research. One may wonder, for instance, if the BFGS updates could be performed at the different grid levels and the resulting matrices prolonged to the fine grid. One may also consider the effect of using simultaneous updates or other quasi-Newton updates in the context of multi-secant algorithms, or the use of the resulting approximations in multilevel optimization algorithms of the type suggested by Nash [25], Gratton, Mouffe, Toint and Weber-Mendonça [13] or Wen and Goldfarb [32]. Yet further research questions include the effect of the particular choice of the restriction and prolongation operators and the impact of grid refinement strategies.

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